

Balancing Communication and Computations in Gradient Tracking Methods for Decentralized Optimization

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Conference on Optimization





Problem

$$\min_{x\in\mathbb{R}^d}f(x)=\sum_{i=1}^n f_i(x)$$

Each function f_i is only known to agent $i \forall i = 1, 2, ..., n$





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Each function f_i is only known to agent $i \forall i = 1, 2, ..., n$





Consensus Optimization Problem

$$egin{aligned} \min_{x_i \in \mathbb{R}^d} \sum_{i=1}^n f_i(x_i) \ s.t. \quad x_i = x_j \quad orall \ i,j \in \mathcal{E} \end{aligned}$$

Each node keeps a local copy $x_i \forall i = 1, 2, ..., n$





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Consensus Optimization Problem

$$\min_{x_i \in \mathbb{R}^d} f(\mathbf{x}) = \sum_{i=1}^n f_i(x_i)$$

s.t. $(\mathbf{W} \otimes I_d)\mathbf{x} = \mathbf{x}$

- x is a concatenation of all local x_i's
- **W** is a symmetric doubly-stochastic matrix that defines the connections in the network

$$\mathbf{X} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \in \mathbb{R}^{nd}, \quad \mathbf{W} = \begin{bmatrix} w_{11} & w_{12} & \cdots & w_{1n} \\ w_{21} & w_{22} & \cdots & w_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ w_{n1} & w_{n2} & \cdots & w_{nn} \end{bmatrix} \in \mathbb{R}^{n \times n}$$



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Consensus Optimization Problem

$$\min_{x_i \in \mathbb{R}^d} f(\mathbf{x}) = \sum_{i=1}^n f_i(x_i)$$

s.t. $\mathbf{Z} \mathbf{x} = \mathbf{x}$

- **x** is a concatenation of all local *x_i*'s
- W is a symmetric doubly-stochastic matrix that defines the connections in the network

$$\mathbf{Z} = \mathbf{W} \otimes I_d \in \mathbb{R}^{nd \times nd}$$

Literature Review

1. Sublinearly Converging Methods

DGD [Bertsekas, Tsitsiklis, et al. 1989, Nedic and Ozdaglar 2009, Sundhar Ram et al. 2010, Tsianos et al. 2012], **NN** [Mokhtari et al. 2017], **NEAR-DGD** [Berahas et al. 2018], ...

2. Linearly Converging Methods

 Push-pull [Pu, Shi, et al. 2020], DlGing [Nedic, Olshevsky, et al. 2017], EXTRA [Shi et al. 2015], SONATA [Sun et al. 2022], NEXT [Di Lorenzo and Scutari 2015],

 Aug-DGM [Xu et al. 2015], LU-GT [Nguyen et al. 2022], ...

3. Asynchronous Methods

[Bertsekas, Tsitsiklis, et al. 1989], [Ram, Veeravalli, and Nedic 2009], **HOGWILD** [Recht et al. 2011], [Wei and Ozdaglar 2013], ...

4. Stochastic Algorithms

DSGT and GSGT [Pu and Nedić 2021], ProxiSkip [Mishchenko et al. 2022], ...

Gradient Tracking Methods

Push-pull [Pu, Shi, et al. 2020], **DIGing** [Nedic, Olshevsky, et al. 2017], **EXTRA** [Shi et al. 2015], **SONATA** [Sun et al. 2022], **NEXT** [Di Lorenzo and Scutari 2015], **Aug-DGM** [Xu et al. 2015], ...

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z} \, \mathbf{x}_k - \alpha \mathbf{y}_k, \\ \mathbf{y}_{k+1} &= \mathbf{Z} \, \mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k) \end{aligned}$$

$$\mathbf{x}_{k} = \begin{bmatrix} x_{1,k} \\ x_{2,k} \\ \vdots \\ x_{n,k} \end{bmatrix} \in \mathbb{R}^{nd}, \quad \mathbf{y}_{k} = \begin{bmatrix} y_{1,k} \\ y_{2,k} \\ \vdots \\ y_{n,k} \end{bmatrix} \in \mathbb{R}^{nd}, \quad \nabla \mathbf{f}(\mathbf{x}_{k}) = \begin{bmatrix} \nabla f_{1}(x_{1,k}) \\ \nabla f_{2}(x_{2,k}) \\ \vdots \\ \nabla f_{n}(x_{n,k}) \end{bmatrix} \in \mathbb{R}^{nd}$$

- Use an additional dual variable \mathbf{y}_k to track the gradient
- Constant α : Linear converge to solution

DIG

Gradient Tracking Methods

Push-pull [Pu, Shi, et al. 2020], **DIGing** [Nedic, Olshevsky, et al. 2017], **EXTRA** [Shi et al. 2015], **SONATA** [Sun et al. 2022], **NEXT** [Di Lorenzo and Scutari 2015], **Aug-DGM** [Xu et al. 2015], ...

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Aug-DGM:
$$\mathbf{x}_{k+1} = \mathbf{Z}(\mathbf{x}_k - \alpha \mathbf{y}_k),$$
 $\mathbf{y}_{k+1} = \mathbf{Z}(\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k))$

Gradient Tracking Methods

Push-pull [Pu, Shi, et al. 2020], **DIGing** [Nedic, Olshevsky, et al. 2017], **EXTRA** [Shi et al. 2015], **SONATA** [Sun et al. 2022], **NEXT** [Di Lorenzo and Scutari 2015], **Aug-DGM** [Xu et al. 2015], ...

DIGing:
$$\mathbf{x}_{k+1} = \mathbf{Z} \mathbf{x}_k - \alpha \mathbf{y}_k,$$

 $\mathbf{y}_{k+1} = \mathbf{Z} \mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)$

Aug-DGM:

$$\begin{array}{l} \mathbf{x}_{k+1} = \mathbf{Z}(\mathbf{x}_k - \alpha \mathbf{y}_k), \\ \mathbf{y}_{k+1} = \mathbf{Z}(\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)) \end{array}$$

 Choice of information shared affects both convergence and practical implementation D

Gradient Tracking Methods

Push-pull [Pu, Shi, et al. 2020], **DIGing** [Nedic, Olshevsky, et al. 2017], **EXTRA** [Shi et al. 2015], **SONATA** [Sun et al. 2022], **NEXT** [Di Lorenzo and Scutari 2015], **Aug-DGM** [Xu et al. 2015], ...

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Aug-DGM:
$$\mathbf{x}_{k+1} = \mathbf{Z}(\mathbf{x}_k - \alpha \mathbf{y}_k),$$
 $\mathbf{y}_{k+1} = \mathbf{Z}(\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k))$

- Choice of information shared affects both convergence and practical implementation
- Applications require a different composition communication and computation steps to achieve overall efficiency



This Talk

- 1. We develop a gradient tracking algorithmic framework (GTA) to unify gradient tracking methods.
- 2. Provide flexibility in number of communication and computation steps in each iteration in GTA.
- 3. Provide sufficient conditions for linear rate of convergence.
- 4. Illustrate benefits of this flexibility with numerical experiments.

GTA Framework

- $\mathbf{W} \in \mathbb{R}^{n \times n} \rightarrow \text{mixing matrix}$
 - Symmetric, Doubly Stochastic
 - Represents the network, i.e., $w_{ii} > 0$ and $w_{ii} > 0$ iff $(i, j) \in \mathcal{E}$

$$\left\| \mathbf{W} - \frac{\mathbf{1}_n \mathbf{1}_n^T}{n} \right\|_2 = \beta \in [0, 1)$$

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 $W_1, W_2, W_3, W_4 \in \mathbb{R}^{n \times n} \rightarrow \text{communication matrices}$

- Symmetric, Doubly Stochastic
- Represents a subset of edges of the network, i.e., $w_{1,ii} > 0$ and $w_{1,ii} \geq 0$ if $(i, j) \in \mathcal{E}$ else $w_{1,ii} = 0$ 11

$$\left\| \mathbf{W}_i - \frac{\mathbf{l}_n \mathbf{l}_n'}{n} \right\|_2 = \beta_i \in [0, 1] \quad \forall \quad i = 1, 2, 3, 4$$

GTA Framework

 $\textbf{W}_1, \textbf{W}_2, \textbf{W}_3, \textbf{W}_4 \rightarrow \text{communication matrices}$

Single communication and computation step in each iteration.

 $\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}_1 \mathbf{x}_k - \alpha \mathbf{Z}_2 \mathbf{y}_k, \\ \mathbf{y}_{k+1} &= \mathbf{Z}_3 \mathbf{y}_k + \mathbf{Z}_4 (\nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)) \end{aligned}$

where $\mathbf{Z}_i = \mathbf{W}_i \otimes I_d \in \mathbb{R}^{nd \times nd} \quad \forall \quad i = 1, 2, 3, 4$

GTA Framework Special Cases

Mixing matrix **W** and $\mathbf{Z} = \mathbf{W} \otimes I_d$

GTA-1 (DIGing, EXTRA, ...)

$$\mathbf{x}_{k+1} = \mathbf{Z}\mathbf{x}_k - \alpha \mathbf{y}_k$$
$$\mathbf{y}_{k+1} = \mathbf{Z}\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)$$

GTA-2 (NEXT, SONATA, ...)

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z} \left(\mathbf{x}_k - \alpha \mathbf{y}_k \right) \\ \mathbf{y}_{k+1} &= \mathbf{Z} \mathbf{y}_k + \nabla \mathbf{f} (\mathbf{x}_{k+1}) - \nabla \mathbf{f} (\mathbf{x}_k) \end{aligned}$$

GTA-3 (Aug-DGM, ATC-DIGing, ...)

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z} \left(\mathbf{x}_k - \alpha \mathbf{y}_k \right) \\ \mathbf{y}_{k+1} &= \mathbf{Z} \left(\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k) \right) \end{aligned}$$



GTA Framework - Convergence Analysis

Definitions

$$\bar{x}_k = \frac{1}{n} \sum_{i=1}^n x_{i,k}, \quad \bar{y}_k = \frac{1}{n} \sum_{i=1}^n y_{i,k}$$



GTA Framework - Convergence Analysis

Definitions

$$\bar{x}_k = \frac{1}{n} \sum_{i=1}^n x_{i,k}, \quad \bar{y}_k = \frac{1}{n} \sum_{i=1}^n y_{i,k}$$

$$r_{k} = \begin{bmatrix} \|\bar{\mathbf{x}}_{k} - \mathbf{x}^{*}\|_{2} \\ \|\mathbf{x}_{k} - \bar{\mathbf{x}}_{k}\|_{2} \\ \|\mathbf{y}_{k} - \bar{\mathbf{y}}_{k}\|_{2} \end{bmatrix}, \quad \bar{\mathbf{x}}_{k} = \begin{bmatrix} \bar{\mathbf{x}}_{k} \\ \bar{\mathbf{x}}_{k} \\ \vdots \\ \bar{\mathbf{x}}_{k} \end{bmatrix} \in \mathbb{R}^{nd}, \quad \bar{\mathbf{y}}_{k} = \begin{bmatrix} \bar{y}_{k} \\ \bar{y}_{k} \\ \vdots \\ \bar{y}_{k} \end{bmatrix} \in \mathbb{R}^{nd}$$



GTA Framework - Convergence Analysis

Definitions

$$\bar{x}_k = \frac{1}{n} \sum_{i=1}^n x_{i,k}, \quad \bar{y}_k = \frac{1}{n} \sum_{i=1}^n y_{i,k}$$

$$r_{k} = \begin{bmatrix} \|\bar{\mathbf{x}}_{k} - \mathbf{x}^{*}\|_{2} \\ \|\mathbf{y}_{k} - \bar{\mathbf{y}}_{k}\|_{2} \end{bmatrix}, \quad \bar{\mathbf{x}}_{k} = \begin{bmatrix} \bar{\mathbf{x}}_{k} \\ \bar{\mathbf{x}}_{k} \\ \vdots \\ \bar{\mathbf{x}}_{k} \end{bmatrix} \in \mathbb{R}^{nd}, \quad \bar{\mathbf{y}}_{k} = \begin{bmatrix} \bar{y}_{k} \\ \bar{y}_{k} \\ \vdots \\ \bar{y}_{k} \end{bmatrix} \in \mathbb{R}^{nd}$$

Assumption

1. The function f is $\mu > 0$ strongly convex and each component function f_i has L > 0 Lipschitz continuous gradients.



GTA Framework - Step size condition

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}_1 \mathbf{x}_k - \alpha \mathbf{Z}_2 \mathbf{y}_k, \\ \mathbf{y}_{k+1} &= \mathbf{Z}_3 \mathbf{y}_k + \mathbf{Z}_4 (\nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)) \end{aligned}$$

Theorem

Suppose Assumption 1 holds and $\beta_1, \beta_3 < 1$ in GTA Framework, then $||r_k||_2$ goes to 0 at a linear rate if

$$\alpha < \min\left\{\frac{1}{L}, \frac{1-\beta_3}{L\beta_4}, \frac{(1-\beta_1+2\beta_2)}{2\beta_2\kappa(L+\mu)}\left(\sqrt{1+\frac{4(1-\beta_1)(1-\beta_3)\beta_2(\kappa+1)}{\beta_4(1-\beta_1+2\beta_2)^2}} - 1\right)\right\}$$
where $\kappa = \frac{L}{\mu}$.



GTA Framework Cases - Step size condition

Theorem

Suppose Assumption 1 holds, $||r_k||_2$ goes to 0 at a linear rate for the special cases if

$$GTA-1: \quad \alpha < \min\left\{\frac{1-\beta}{L}, \frac{(3-\beta)}{2\kappa(L+\mu)}\left(\sqrt{1+4(\kappa+1)\left(\frac{1-\beta}{3-\beta}\right)^2}-1\right)\right\}$$
$$GTA-2: \quad \alpha < \min\left\{\frac{1-\beta}{L}, \frac{(1+\beta)}{2\kappa(L+\mu)\beta}\left(\sqrt{1+4(\kappa+1)\beta\left(\frac{1-\beta}{1+\beta}\right)^2}-1\right)\right\}$$

GTA-3:
$$\alpha < \min\left\{\frac{1}{L}, \frac{1-\beta}{L\beta}, \frac{(1+\beta)}{2\kappa(L+\mu)\beta}\left(\sqrt{1+4(\kappa+1)\left(\frac{1-\beta}{1+\beta}\right)^2}-1\right)\right\}$$

where $\kappa = \frac{L}{\mu}$.



GTA Framework Cases - Rate of Convergence

Theorem

Suppose Assumption 1 holds and $\alpha \leq \frac{1}{L} ||r_k||_2$ goes to 0 at a linear rate upper bounded by the following expressions

$$GTA-1: \max\left\{1 - \frac{\alpha\mu}{2}, \ \beta + \sqrt{\alpha L}\left(2.5 + \sqrt{\kappa}\right)\right\}$$
$$GTA-2: \max\left\{1 - \frac{\alpha\mu}{2}, \ \beta + \sqrt{\alpha L}\left(2.5 + \sqrt{\kappa\beta}\right)\right\}$$
$$GTA-3: \max\left\{1 - \frac{\alpha\mu}{2}, \ \beta\left(1 + \sqrt{\alpha L}\left(2.5 + \sqrt{\kappa}\right)\right)\right\}$$

where $\kappa = \frac{L}{\mu}$.

GTA Framework - Numerical Experiments

Almost Full network $\beta = 0.25$



Figure: Quadratics, n = 16, d = 10, $\kappa = 10^4$

GTA Framework - Numerical Experiments

Cyclic Network $\beta = 0.992$



Figure: Quadratics, n = 16, d = 10, $\kappa = 10^4$



(a) Single Communication

$$\mathbf{W} = \begin{bmatrix} 0.9 & 0.1 & 0 & 0 \\ 0.1 & 0.8 & 0.1 & 0 \\ 0 & 0.1 & 0.8 & 0.1 \\ 0 & 0 & 0.1 & 0.9 \end{bmatrix}$$



(a) Single Communication



(b) 2 Communications

$$\mathbf{W} = \begin{bmatrix} 0.9 & 0.1 & 0 & 0 \\ 0.1 & 0.8 & 0.1 & 0 \\ 0 & 0.1 & 0.8 & 0.1 \\ 0 & 0 & 0.1 & 0.9 \end{bmatrix}$$



(a) Single Communication



(b) 2 Communications

$$\mathbf{W} = \begin{bmatrix} 0.9 & 0.1 & 0 & 0 \\ 0.1 & 0.8 & 0.1 & 0 \\ 0 & 0.1 & 0.8 & 0.1 \\ 0 & 0 & 0.1 & 0.9 \end{bmatrix} \qquad \mathbf{W}^2 = \begin{bmatrix} 0.82 & 0.17 & 0.01 & 0 \\ 0.17 & 0.66 & 0.16 & 0.01 \\ 0.01 & 0.16 & 0.66 & 0.17 \\ 0 & 0.01 & 0.17 & 0.82 \end{bmatrix}$$



(a) Single Communication



(b) 200 Communications



(a) Single Communication



(b) 200 Communications

$$\mathbf{W} = \begin{bmatrix} 0.9 & 0.1 & 0 & 0 \\ 0.1 & 0.8 & 0.1 & 0 \\ 0 & 0.1 & 0.8 & 0.1 \\ 0 & 0 & 0.1 & 0.9 \end{bmatrix} \qquad \mathbf{W}^{200} = \begin{bmatrix} 0.25 & 0.25 & 0.25 & 0.25 \\ 0.25 & 0.25 & 0.25 & 0.25 \\ 0.25 & 0.25 & 0.25 & 0.25 \\ 0.25 & 0.25 & 0.25 & 0.25 \end{bmatrix}$$



GTA Framework - Multiple Communications

 $n_c \rightarrow \#$ of communication steps $W_i \rightarrow W_i^{n_c}$ $\forall i = 1, 2, 3, 4$ $\beta_i \rightarrow \beta_i^{n_c}$ $\forall i = 1, 2, 3, 4$ $\mathbf{Z}_i \rightarrow \mathbf{Z}_i^{\mathbf{n_c}} = \mathbf{W}_i^{\mathbf{n_c}} \otimes I_d$ $\forall i = 1, 2, 3, 4$



GTA Framework - Multiple Communications

 $n_c \rightarrow \#$ of communication steps $W_i \rightarrow W_i^{n_c}$ $\forall i = 1, 2, 3, 4$ $\beta_i \rightarrow \beta_i^{n_c}$ $\forall i = 1, 2, 3, 4$ $\mathbf{Z}_i \rightarrow \mathbf{Z}_i^{\mathbf{n_c}} = \mathbf{W}_i^{\mathbf{n_c}} \otimes I_d$ $\forall i = 1, 2, 3, 4$

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}_1^{n_c} \mathbf{x}_k - \alpha \mathbf{Z}_2^{n_c} \mathbf{y}_k, \\ \mathbf{y}_{k+1} &= \mathbf{Z}_3^{n_c} \mathbf{y}_k + \mathbf{Z}_4^{n_c} (\nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)) \end{aligned}$$



GTA Framework - Multiple Communications

 $\begin{array}{l} \textbf{n_c} \to \# \text{of communication steps} \\ \textbf{W}_i \to \textbf{W}_i^{n_c} & \forall i = 1, 2, 3, 4 \\ \beta_i \to \beta_i^{n_c} & \forall i = 1, 2, 3, 4 \\ \textbf{Z}_i \to \textbf{Z}_i^{n_c} = \textbf{W}_i^{n_c} \otimes I_d & \forall i = 1, 2, 3, 4 \end{array}$

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}_1^{n_c} \mathbf{x}_k - \alpha \mathbf{Z}_2^{n_c} \mathbf{y}_k, \\ \mathbf{y}_{k+1} &= \mathbf{Z}_3^{n_c} \mathbf{y}_k + \mathbf{Z}_4^{n_c} (\nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k)) \end{aligned}$$

With more communcation, i.e., increase in nc

- The step size condition increases
- The rate of convergence decreases

GTA Special Cases - Multiple Communications

Mixing matrix **W** and $\mathbf{Z}^{n_c} = \mathbf{W}^{n_c} \otimes I_d$

GTA-1

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}^{n_c} \mathbf{x}_k - \alpha \mathbf{y}_k \\ \mathbf{y}_{k+1} &= \mathbf{Z}^{n_c} \mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k) \end{aligned}$$

GTA-2

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}^{n_c} \left(\mathbf{x}_k - \alpha \mathbf{y}_k \right) \\ \mathbf{y}_{k+1} &= \mathbf{Z}^{n_c} \mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k) \end{aligned}$$

GTA-3

$$\begin{aligned} \mathbf{x}_{k+1} &= \mathbf{Z}^{n_c} \left(\mathbf{x}_k - \alpha \mathbf{y}_k \right) \\ \mathbf{y}_{k+1} &= \mathbf{Z}^{n_c} \left(\mathbf{y}_k + \nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k) \right) \end{aligned}$$



GTA Multiple Communications - Rate of Convergence

Theorem

Suppose Assumption 1 holds, number of communications is at least 1 $(n_c \ge 1)$ and $\alpha \le \frac{1}{L}$, $||r_k||_2$ goes to 0 at a linear rate upper bounded by the following expressions

For GTA-1
$$\max\left\{1-\frac{\alpha\mu}{2}, \ \beta^{n_c}+\sqrt{\alpha L}\left(2.5+\sqrt{\kappa}\right)\right\}$$

For GTA-2 $\max\left\{1-\frac{\alpha\mu}{2}, \ \beta^{n_c}+\sqrt{\alpha L}\left(2.5+\sqrt{\kappa\beta^{n_c}}\right)\right\}$
For GTA-3 $\max\left\{1-\frac{\alpha\mu}{2}, \ \beta^{n_c}\left(1+\sqrt{\alpha L}\left(2.5+\sqrt{\kappa}\right)\right)\right\}$

where $\kappa = \frac{L}{\mu}$.

GTA Multiple Communications - Numerical Experiments



Figure: Quadratics, n = 16, d = 10, $\kappa = 10^4$

GTA Multiple Communications - Numerical Experiments

Cyclic Network $\beta = 0.992$ 10^{-1} 10^{1} 10^{-2} Optimization Error 10⁻¹ Consensus Error 10^{-3} 10-4 10^{-5} 10^{-2} 10^{-6} 4 1e4 Communications Communications 1e4 GTA-1 (1, 1) GTA-2 (1, 1) GTA-3 (1, 1) GTA-1 (5, 1) ---- GTA-2 (5, 1) ------ GTA-3 (5, 1) ---¥----

Figure: Quadratics, n = 16, d = 10, $\kappa = 10^4$



 $n_c \rightarrow \#$ of communication steps $n_g \rightarrow \#$ of computation steps

$$\begin{split} \mathbf{x}_{k+1,1} &= \mathbf{Z}_1^{n_c} \mathbf{x}_{k,n_g} - \alpha \mathbf{Z}_2^{n_c} \mathbf{y}_{k,n_g} \\ \mathbf{y}_{k+1,1} &= \mathbf{Z}_3^{n_c} \mathbf{y}_{k,n_g} + \mathbf{Z}_4^{n_c} (\nabla \mathbf{f}(\mathbf{x}_{k+1,1}) - \nabla \mathbf{f}(\mathbf{x}_{k,n_g})) \end{split}$$



 $n_c \rightarrow \#$ of communication steps $n_g \rightarrow \#$ of computation steps

$$\begin{aligned} \mathbf{x}_{k+1,1} &= \mathbf{Z}_{1}^{n_{c}} \mathbf{x}_{k,n_{g}} - \alpha \mathbf{Z}_{2}^{n_{c}} \mathbf{y}_{k,n_{g}} \\ \mathbf{y}_{k+1,1} &= \mathbf{Z}_{3}^{n_{c}} \mathbf{y}_{k,n_{g}} + \mathbf{Z}_{4}^{n_{c}} (\nabla \mathbf{f}(\mathbf{x}_{k+1,1}) - \nabla \mathbf{f}(\mathbf{x}_{k,n_{g}})) \\ \text{For } j \rightarrow 1, 2, ..., n_{g} - 1 \\ \mathbf{x}_{k+1,j+1} &= \mathbf{x}_{k+1,j} - \alpha \mathbf{y}_{k+1,j}, \\ \mathbf{y}_{k+1,j+1} &= \mathbf{y}_{k+1,j+1} + \nabla \mathbf{f}(\mathbf{x}_{k+1,j+1}) - \nabla \mathbf{f}(\mathbf{x}_{k+1,j}) \end{aligned}$$



 $n_c \rightarrow \#$ of communication steps $n_g \rightarrow \#$ of computation steps

Theorem

Under previous assumptions, $\beta_1, \beta_3 < 1$, number of communication steps is at least one ($n_c \ge 1$) and number of computation steps is finite ($1 \le n_g < \infty$), then $\exists \alpha > 0$, s.t. $||r_k||_2$ goes to 0 at a linear rate.



 $n_c \rightarrow \#$ of communication steps $n_g \rightarrow \#$ of computation steps

Theorem

Under previous assumptions, $\beta_1, \beta_3 < 1$, number of communication steps is at least one ($n_c \ge 1$) and number of computation steps is finite ($1 \le n_g < \infty$), then $\exists \alpha > 0$, s.t. $||r_k||_2$ goes to 0 at a linear rate.

- The step size increases with an increase in n_c, i.e., number of communication steps.
- The step size is inversely proportional to n_g, i.e., number of computation steps.



GTA Special Cases

$$GTA-1 \qquad \mathbf{x}_{k+1,1} = \mathbf{Z}^{n_c} \mathbf{x}_{k,n_g} - \alpha \mathbf{y}_{k,n_g}$$
$$\mathbf{y}_{k+1,1} = \mathbf{Z}^{n_c} \mathbf{y}_{k,n_g} + \nabla \mathbf{f}(\mathbf{x}_{k+1,1}) - \nabla \mathbf{f}(\mathbf{x}_{k,n_g})$$
$$\rightarrow n_g - 1 \text{ compute steps}$$

GTA-2

$$\mathbf{x}_{k+1,1} = \mathbf{Z}^{n_c} \left(\mathbf{x}_{k,n_g} - \alpha \mathbf{y}_{k,n_g} \right)$$

$$\mathbf{y}_{k+1,1} = \mathbf{Z}^{n_c} \mathbf{y}_{k,n_g} + \nabla \mathbf{f}(\mathbf{x}_{k+1,1}) - \nabla \mathbf{f}(\mathbf{x}_{k,n_g})$$

$$\rightarrow n_g - 1 \text{ compute steps}$$

GTA-3

$$\mathbf{x}_{k+1,1} = \mathbf{Z}^{n_c} \left(\mathbf{x}_{k,n_g} - \alpha \mathbf{y}_{k,n_g} \right)$$
$$\mathbf{y}_{k+1,1} = \mathbf{Z}^{n_c} \left(\mathbf{y}_{k,n_g} + \nabla \mathbf{f}(\mathbf{x}_{k+1,1}) - \nabla \mathbf{f}(\mathbf{x}_{k,n_g}) \right)$$
$$\rightarrow n_g - 1 \text{ compute steps}$$



GTA Special Cases

 $n_c \rightarrow \#$ of communication steps $n_g \rightarrow \#$ of computation steps

If the same step size is employed in all three methods, their convergence rates can be ordered as:

 $GTA-3(n_c, n_g) \leq GTA-2(n_c, n_g) \leq GTA-1(n_c, n_g)$

GTA Multiple Communications and Computations -Numerical Experiments



Figure: Quadratics, n = 16, d = 10, $\kappa = 10^4$

GTA Multiple Communications and Computations -Numerical Experiments



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Flexible Randomized Gradient Tracking Algorithm

Performing communications less often randomly FedAvg [McMahan et al. 2017], FedLin [Mitra et al. 2021], Scaffold [Karimireddy et al. 2020], Scaffnew [Mishchenko et al. 2022], ...

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With probability *p*:
$$\mathbf{x}_{k+1} = \mathbf{Z}_1^{n_c} \mathbf{x}_k - \alpha \mathbf{Z}_2^{n_c} \mathbf{y}_k,$$

 $\mathbf{y}_{k+1} = \mathbf{Z}_3^{n_c} \mathbf{y}_k + \mathbf{Z}_4^{n_c} (\nabla \mathbf{f}(\mathbf{x}_{k+1}) - \nabla \mathbf{f}(\mathbf{x}_k))$
Else: $\mathbf{x}_{k+1} = \mathbf{x}_k - \alpha \mathbf{y}_k,$
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- Less rigid, achieves the desired balance in expectation
- Good theoretical and empirical performance
- Paper soon to follow !!



Conclusions

- We provide a unifying gradient tracking algorithmic framework that allows performing theoretical comparisons between different gradient tracking methods.
- 2. We provide the flexibility to perform any composition of communication and computation steps in each iteration and show linear rate of convergence.
- **3**. Adapting your algorithm to the system with this flexibility can allow you to improve convergence rate.

Paper available at: https://arxiv.org/abs/2303.14289



Thank You! Questions?



Backup Slides



$$r_{k+1} \leq A(n_c)r_k$$

$$A(n_c) = \begin{bmatrix} 1 - \alpha \mu & \frac{\alpha L}{\sqrt{n}} & 0\\ 0 & \beta_1^{n_c} & \alpha \beta_2^{n_c}\\ \sqrt{n} \alpha \beta_4^{n_c} L^2 & \beta_4^{n_c} L(\|\mathbf{Z}_1^{n_c} - I_{nd}\|_2 + \alpha L) & \beta_3^{n_c} + \alpha \beta_4^{n_c} L \end{bmatrix}$$



$$\begin{aligned} r_{k+1} &\leq B(n_c, n_g) r_k, \\ \text{where } B(n_c, n_g) &= A(n_c, n_g) + \alpha L(n_g - 1) E(n_c, n_g) \end{aligned}$$
$$A(n_c, n_g) &= \begin{bmatrix} (1 - \alpha \mu)^{n_g} & \frac{\kappa}{\sqrt{n}} (1 - (1 - \alpha \mu)^{n_g}) & 0 \\ 0 & \beta_1^{n_c} & \alpha ((n_g - 1)\beta_1^{n_c} + \beta_2^{n_c}) \\ \sqrt{n} \alpha \beta_4^{n_c} L^2 & \beta_4^{n_c} L(\|\mathbf{Z}_1^{n_c} - I_{nd}\|_2 + \alpha L) & \beta_3^{n_c} + \alpha \beta_4^{n_c} L \end{bmatrix}$$
$$E(n_c, n_g) &= \begin{bmatrix} \alpha L n_g & \frac{\alpha L n_g}{\sqrt{n}} & \frac{\alpha n_g}{\sqrt{n}} \\ \sqrt{n} \alpha L \delta_1(n_c, n_g) & \alpha L \delta_1(n_c, n_g) & \alpha \delta_1(n_c, n_g) \\ \sqrt{n} L \delta_2(n_c, n_g) & L \delta_2(n_c, n_g) & \delta_2(n_c, n_g) \end{bmatrix}$$

and

$$\begin{split} \delta_1(n_c, n_g) &= 2\beta_2^{n_c} + \beta_1^{n_c}(n_g - 2), \\ \delta_2(n_c, n_g) &= 2\left(\beta_4^{n_c} \|\mathbf{Z}_1^{n_c} - I_{nd}\|_2 + \frac{\beta_4^{n_c}}{n_g} + \beta_3^{n_c}\right). \end{split}$$